Progressive Join Algorithms Considering User Preference




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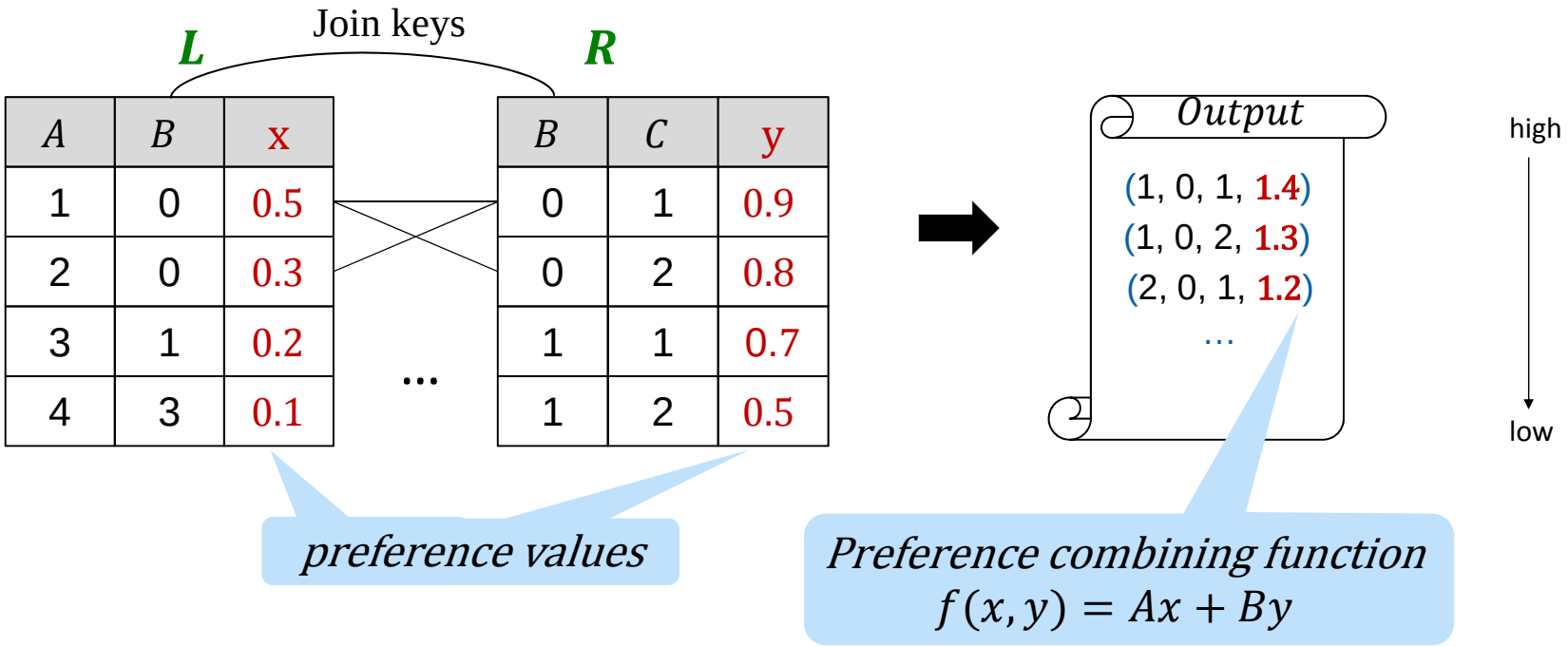
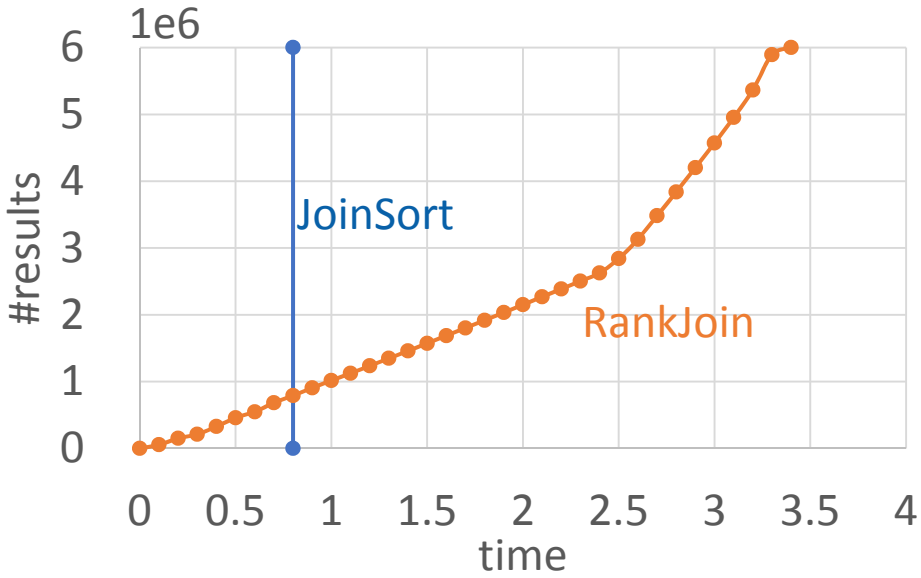
Preference-aware Progressive Join

- Progressive query processing for exploratory data analysis
- Return join results ordered according to preference



- Goal: Fast early results & Fast full results

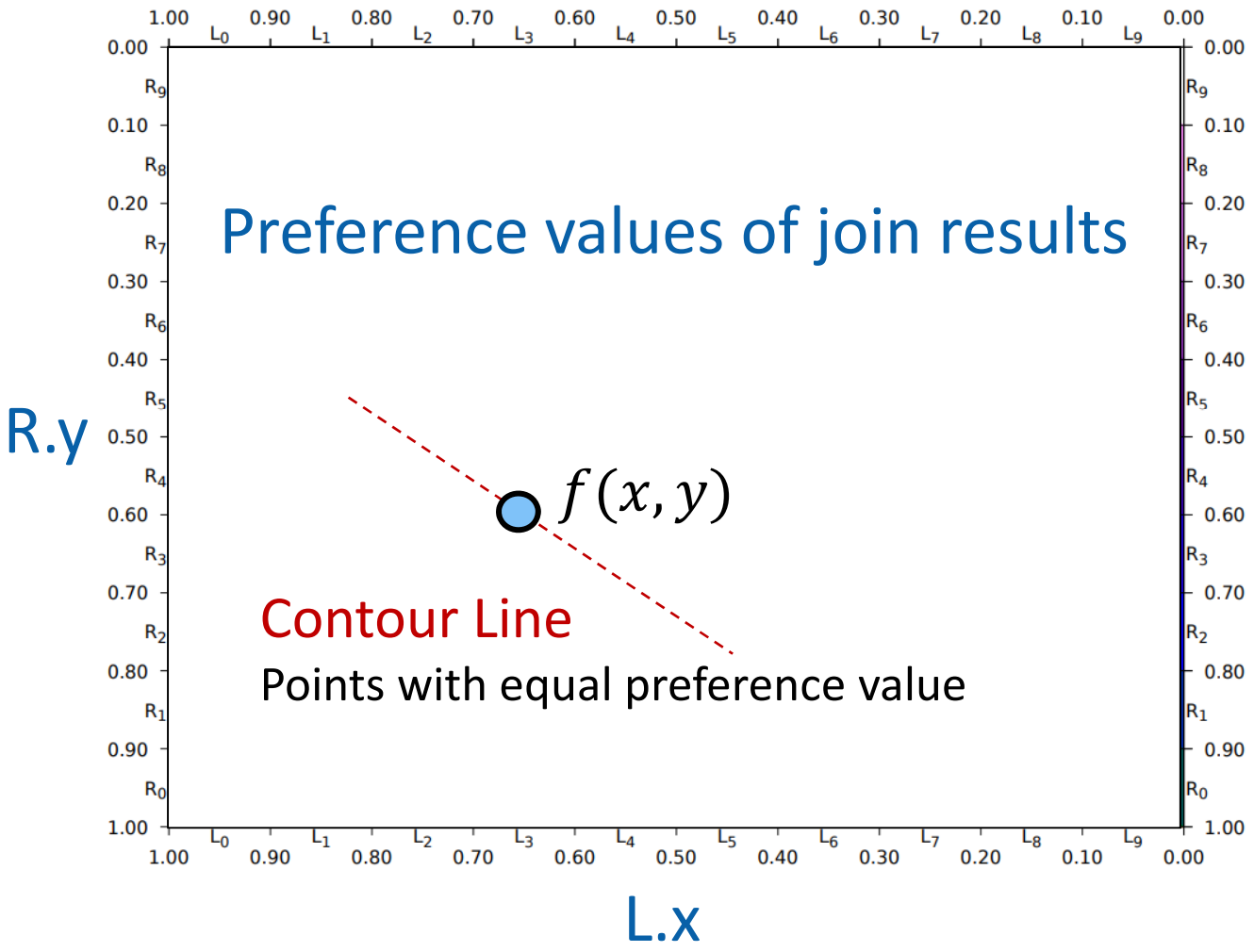
Problems of Existing Solutions

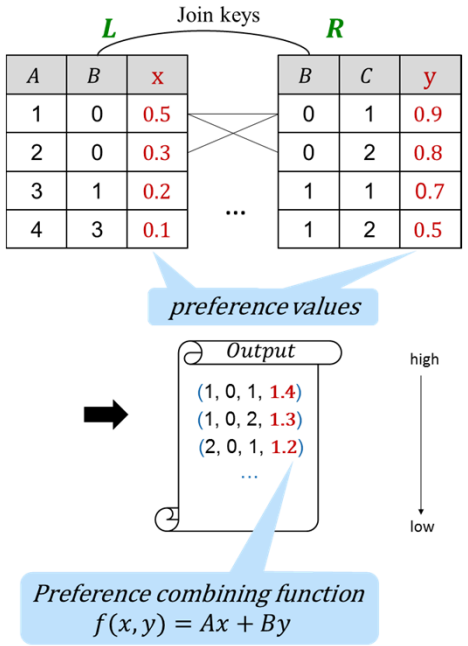


- Blocking approach: Join + Sort
  - **No early results**
- Extending top-k join algorithms:
  - e.g., RankJoin [VLDB'03] symmetric hash join + priority queue
  - **Slow full results**
  - **Significant sorting overhead**

We want to reduce or eliminate sorting overhead

Our Idea: Exploiting Contour Lines





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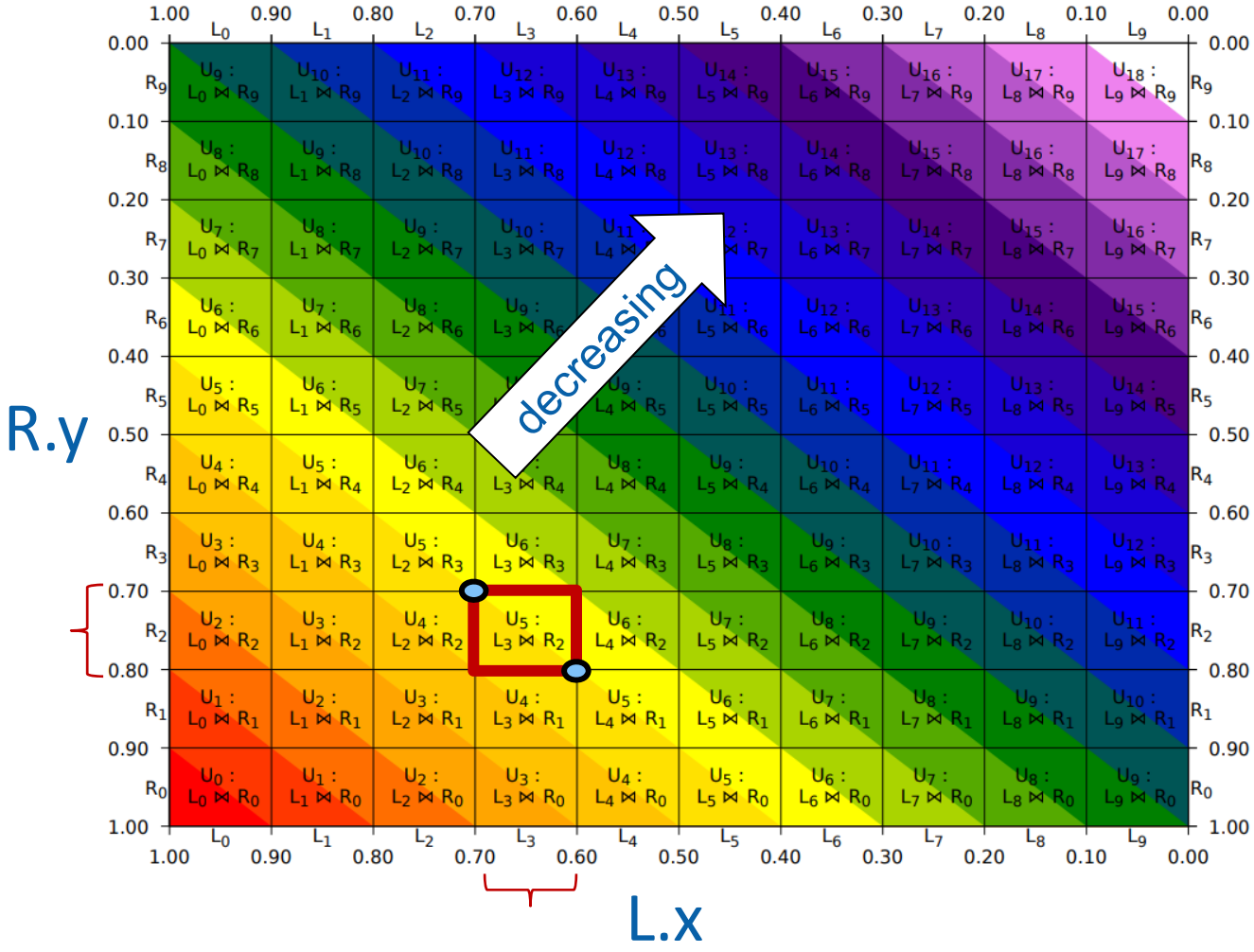
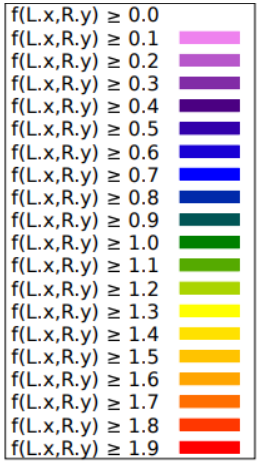
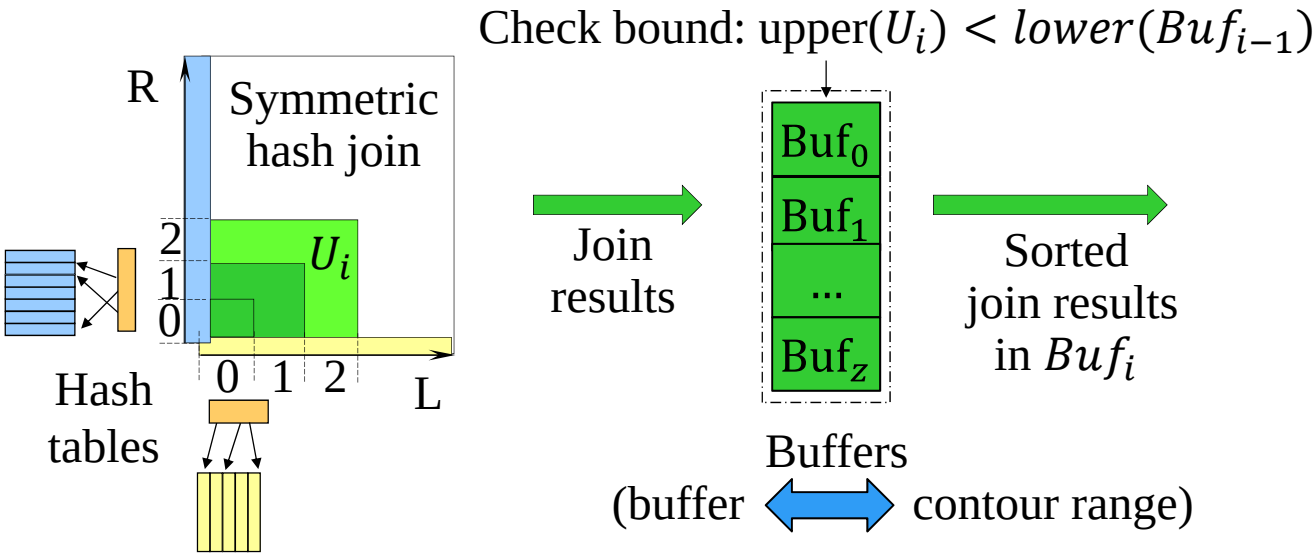




$$f(x, y) = Ax + By$$

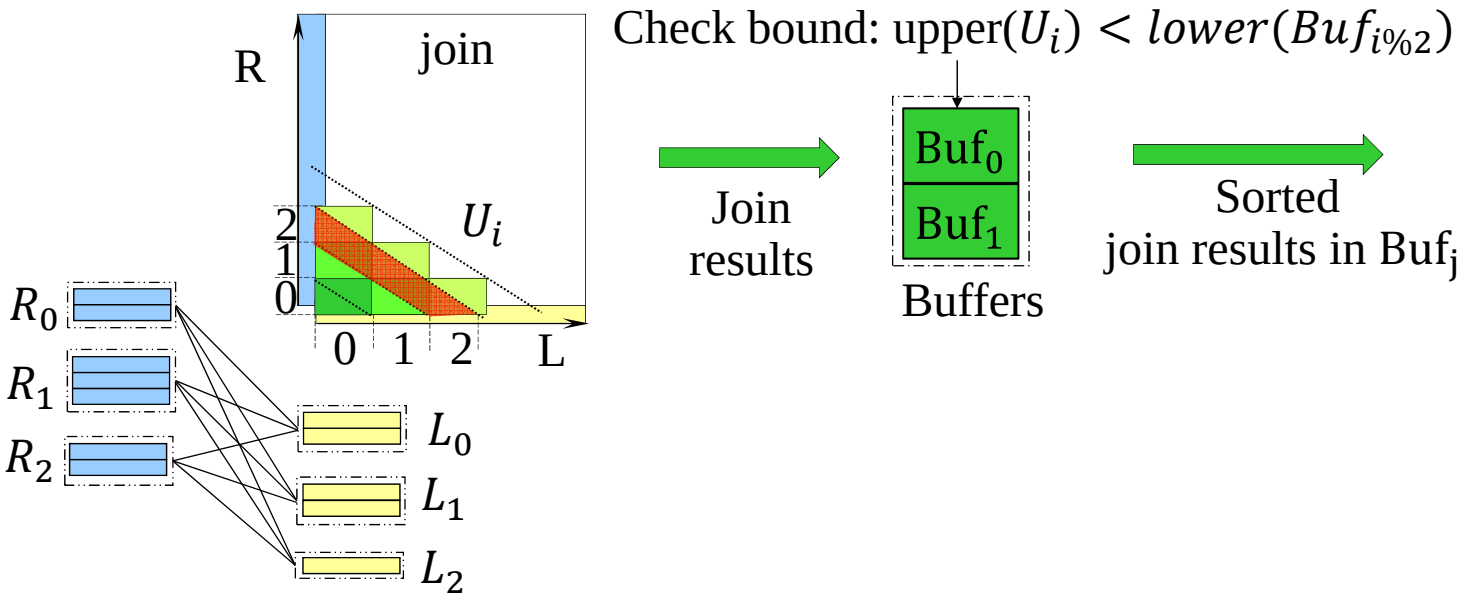
A=B=1

Inputs Follow Contour Lines (Algorithm1: CJpI)



- Avoid sorting across buffers
- Only need to sort within each buffer

Both Inputs and Outputs Follow Contour Lines (Algorithm2: CJpB)



- Reduce intermediate result size
- But join cost may be higher

Relaxation to Remove Sorting

- Relaxation of Order: **no intra-buffer sorting**
  - Relaxation: tuples t_i and t_j are regarded as in order if $\left|t_i.pval - t_j.pval\right| \leq \epsilon$
  - Judiciously select input intervals

4 Variants of Contour Joins

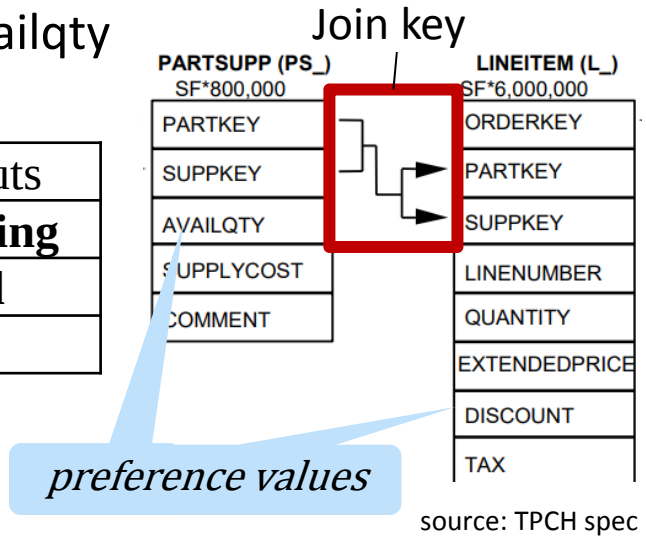
Variants	Follow Contour Lines		Relaxation
	Join Inputs	Join Results	
CJpI	✓		
CJpB	✓	✓	
CJrI	✓		✓
CJrB	✓	✓	✓

p: precise, r: relaxed; I: Inputs, B: Both inputs & outputs

Experiments

- Data Set: Based on TPC-H Lineitem and Partsupp
  - ❑ Preference values based on: l_discount, ps_availqty
  - ❑ 3 datasets (can fit into main memory)

Scale Factor	Datasets	#Inputs	#Outputs
m_1	$datasets\#1$	**fixed**	**increasing**
m_2	$datasets\#2$	**increasing**	**fixed**
m_3	$datasets\#3$	**increasing**	



source: TPCH spec

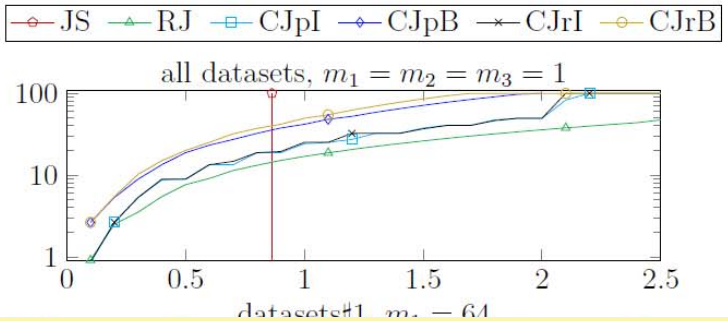
- Query

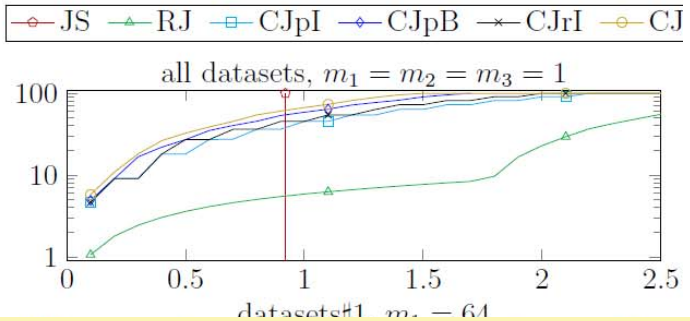
$f(x,y) = Ax + By$

```
select  L.key, f(L.pval, PS.pval) as score
from    Lineitem as L, Partsupp as PS
where   L.key = PS.key
order by score ASC
progressively
```

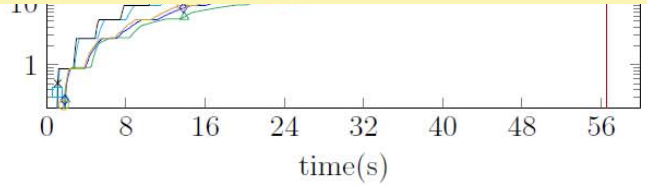
- Machine
  - ❑ Intel Core i7-4770 CPU @3.40GHz (8MB cache) and 32GB memory
  - ❑ 64-bit Ubuntu 16.04 LTS with 4.15.0-62-generic Linux kernel

Overall Results

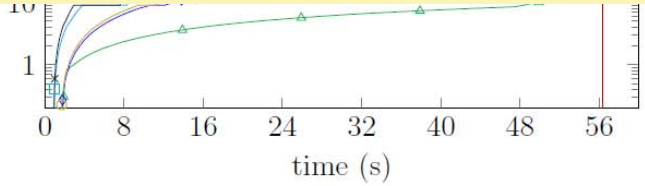




- **Compared to RJ (RankJoin)**
  - 1% early results
    - best **precise** contour join: up to 7x improvements
    - best **relaxed** contour join: up to 14x improvements
  - Full results
    - best **precise** contour join: up to 10.6x improvements
    - best **relaxed** contour join: up to 39.4x improvements
- **Comparable or better perfomance to JS (blocking JoinSort)**

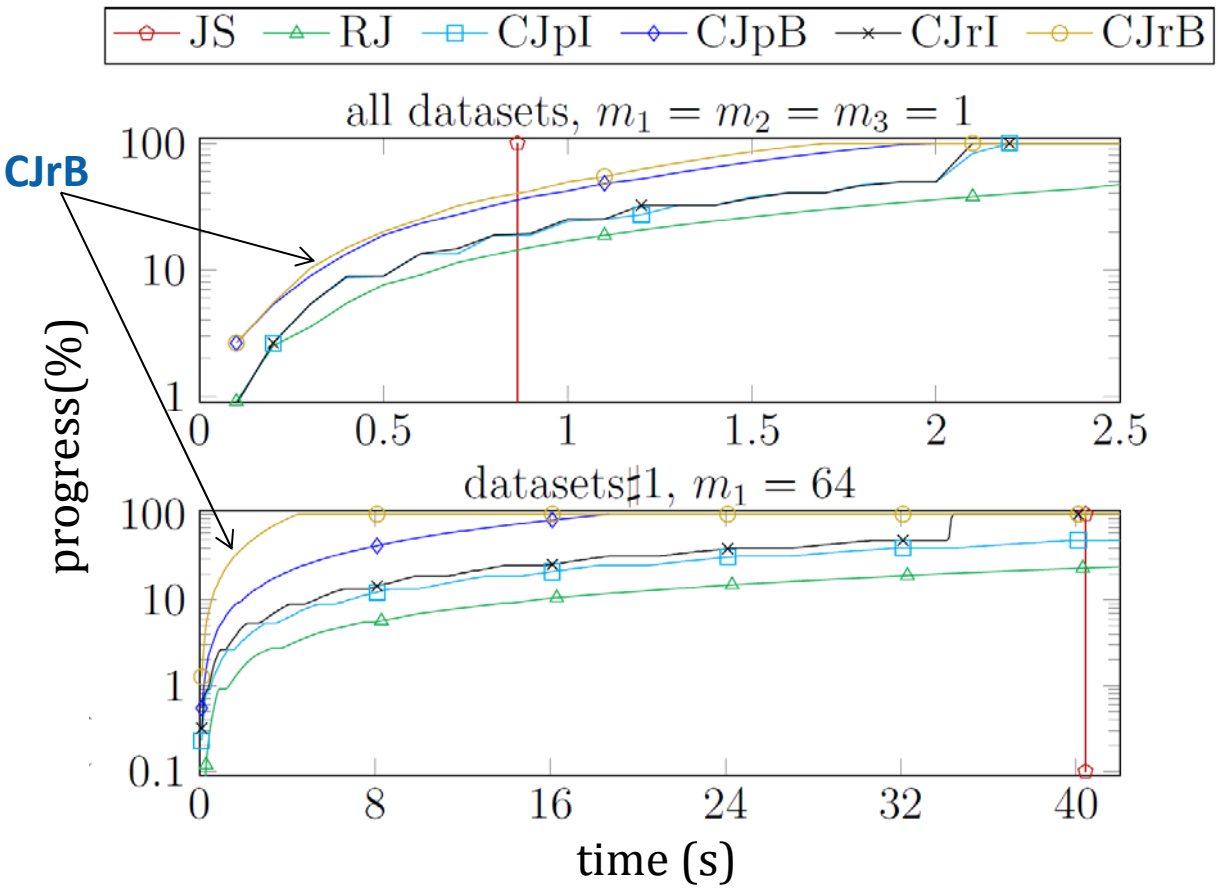


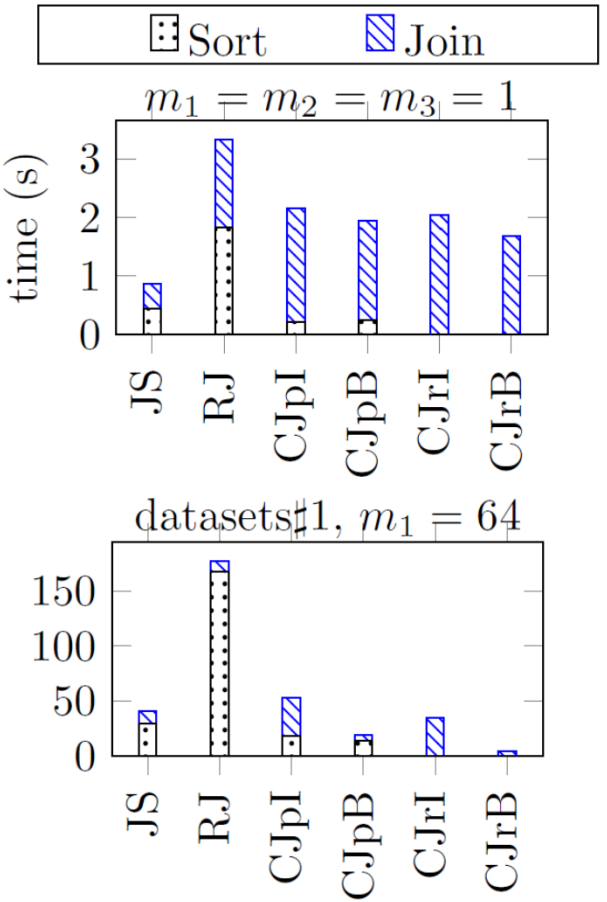
(a) $f(x, y) = x + y$, $p_L = 200$, $p_R = 200$



(b) $f(x, y) = 10x + y$, $p_L = 2000$, $p_R = 200$

Fix Input Size, Increase Join Result Size





- As join result size increases, the fraction of sorting increases
- RankJoin becomes very poor
- CJrB is the best performing algorithm

Conclusion

- Preference-aware joins in progressive query processing
- Idea: exploit contour lines in the join result space
- ContourJoin: a promising solution
  - Faster early and full results generation (vs. RankJoin)
  - Good total result generation performance (vs. JoinSort)

- More in the paper
  - Algorithms and proofs
  - Extensive experimental results
  - Discussion on preference combining functions, unsorted inputs, multi-way joins